## Explicit construction and parameters of projective toric codes

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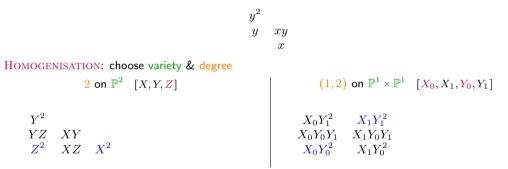
https://arxiv.org/abs/2003.10357

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 Example of classical/Projective toric code

**Classical toric code:** Span of the evaluation on  $(\mathbb{F}_q^*)^2$  of monomials



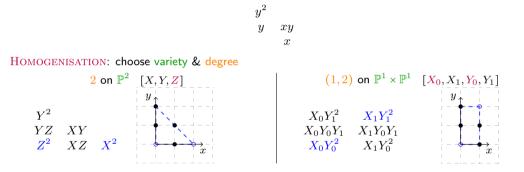
Projective toric code: Span of the evaluation of monomials on rational points of the whole variety

 $\begin{array}{c} (a,b,1) \ (a,1,0) \ (1,0,0) \\ (a,b) \in \mathbb{F}_q^2 \end{array} \begin{array}{c} (1,a,1,b) \ (0,1,1,b) \\ (1,a,0,1) \ (0,1,0,1) \end{array}$ 

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**Classical toric code:** Span of the evaluation on  $(\mathbb{F}_q^*)^2$  of monomials



**Projective toric code:** Span of the evaluation of monomials on rational points of the *whole* variety (1, a, 1, b) (0, 1, 1, b)

(a, b, 1) (a, 1, 0) (1, 0, 0)  $(a, b) \in \mathbb{F}_q^2$ Polygon  $\leftrightarrow$  variety & degree

Explicit construction and parameters of projective toric codes

Introduction 00 Classical/Projective toric codes An integral polytope  $P \subset \mathbb{R}^N$  (vertices in  $\mathbb{Z}^N$ ) defines an abstract toric variety  $\mathbf{X}_P$  with a divisor D and a monomial basis of L(D) (set of polynomials of a certain *degree*). Size of  $P \leftrightarrow \text{Degree}$  in L(D) $\mathbb{D}^1 \setminus \mathbb{P}^1$  $\mathbb{P}^1 \times \mathbb{P}^1 \times \mathbb{P}^1$ Degree (1,2)Degree 2 Degree (4, 3, 3)Why toric?  $X_P$  contains a dense torus  $\mathbb{T}_P \simeq \left(\overline{\mathbb{F}_q}^*\right)^N$  whose rational points are  $(\mathbb{F}_q^*)^N$ . Classical toric code:  $C_P = \{(f(t))_{t \in \mathbb{T}_P}(\mathbb{F}_q) \mid f \in L(D)\}$ Hansen [Han02], Little-Schwarz [LS05], Ruano [Rua07], Soprunov-Soprunova [SS09] Aim : Constructing and studying the projective toric code

 $\mathsf{PC}_P = \left\{ (f(\mathbf{x}))_{\mathbf{x} \in \mathbf{X}_P(\mathbb{F}_q)} \mid f \in L(D) \right\}$ 

### Advantages similar to $RM \rightarrow PRM$ :

 $\bullet$  length  $\nearrow$ , minimum distance  $\nearrow$  with roughly the same dimension.

Strenghten the geometric interpretation

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	Handling a toric variety					
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Description o	f the toric variety ${f X}$	$_P$ associated	l to the polyt	cope P		

Several ways to describe  $\mathbf{X}_P$  thanks to the integral polytope P: (under some assumptions)

- $\oplus$  geometric properties
- with *fans* as an abstract variety
- $\ominus$  implementation

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- as a quotient of a subset of  $\mathbb{A}^r$  (where  $r = \mathsf{nb}$  of facets of P) by a group G (simplicial variety)
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  - $\oplus$  functions of L(D) = polynomials in r variables

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*Example:*  $P = \operatorname{Conv}((0,0), (1,0), (0,1), (1,1)) \subset \mathbb{R}^2$  gives  $\mathbf{X}_P = \mathbb{P}^1 \times \mathbb{P}^1$ :

- embedded in  $\mathbb{P}^3$  by the Segre map:  $(x_0, x_1, y_0, y_1) \mapsto (x_i y_j)$ ,
- defined as the quotient of  $(\mathbb{A}^2 \smallsetminus \{(0,0)\})^2 \subset \mathbb{A}^4$  by the group  $(\overline{\mathbb{F}}^*)^2$  via the action

$$(\lambda,\mu)\cdot(x_0,x_1,y_0,y_1)=(\lambda x_0,\lambda x_1,\mu y_0,\mu y_1)$$

Functions= bihomogeneous polynomials

	Handling a toric variety			
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For classical toric codes, an integral point  $m \in P \cap \mathbb{Z}^N$  gives a monomial  $\chi^m = X_1^{m_1} \dots X_N^{m_N}$ . In the projective case, it corresponds to a monomial  $\chi^{(m,P)} \in \mathbb{F}_q[\mathbf{X}_1, \dots, \mathbf{X}_r]$ .

$$L(D) = \operatorname{Span}\left(\chi^{\langle m, P \rangle} \mid m \in P \cap \mathbb{Z}^N\right)$$

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We can go from  $\chi^m$  to  $\chi^{(m,P)}$  via homogenization process. Example on  $\mathbb{P}^2$ :



- $\chi^m = x_1^0 x_2^1 = x_2.$
- $\chi^{\langle m, P \rangle} = X_2 \leftarrow \text{homogenized in degree } 1$
- $\chi^{\langle m, 2P \rangle} = X_0 X_2 \leftarrow \text{homogenized in degree } 2$

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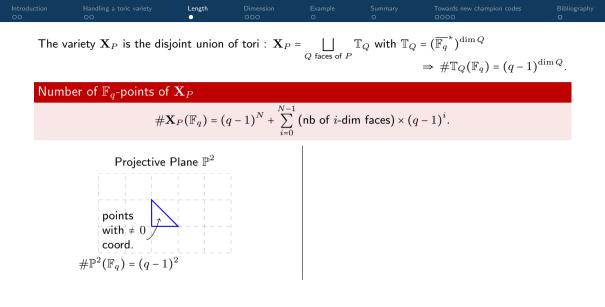
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• 
$$\chi^{(m,2P)} = X_0 X_2 \leftarrow \text{homogenized in degree } 2$$

$$\mathsf{PC}_P = \mathrm{Span}\left\{\left(\chi^{(m,P)}(\mathbf{x})\right)_{\mathbf{x}\in\mathcal{P}}\in\mathbb{F}_q^n, \ m\in P\cap\mathbb{Z}^N\right\} \text{ where } n = \#\mathbf{X}_P(\mathbb{F}_q).$$



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The variety 
$$\mathbf{X}_P$$
 is the disjoint union of tori :  $\mathbf{X}_P = \bigcup_{Q \text{ faces of } P} \mathbb{T}_Q$  with  $\mathbb{T}_Q = (\overline{\mathbb{F}_q}^*)^{\dim Q}$   
 $\Rightarrow \#\mathbb{T}_Q(\mathbb{F}_q) = (q-1)^{\dim Q}$ .  
Number of  $\mathbb{F}_q$ -points of  $\mathbf{X}_P$   
 $\#\mathbf{X}_P(\mathbb{F}_q) = (q-1)^N + \sum_{i=0}^{N-1} (\text{nb of } i\text{-dim faces}) \times (q-1)^i$ .  
Projective Plane  $\mathbb{P}^2$   
 $\#\mathbb{P}^2(\mathbb{F}_q) = (q-1)^2 + 3(q-1)$ 

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$$\frac{\text{Production}}{\text{OC}} \xrightarrow{\text{Parading a toric variety}} \underbrace{\text{Length}}_{\text{OC}} \xrightarrow{\text{Dimension}}_{\text{OC}} \underbrace{\text{Summary}}_{\text{OC}} \xrightarrow{\text{Deverde a revert champion code}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \underbrace{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \underbrace{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \underbrace{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \underbrace{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \underbrace{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}} \xrightarrow{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}} \xrightarrow{\text{OCO}}_{\text{OCO}} \xrightarrow{\text{OCO}} \xrightarrow{\text{OCO$$

		Dimension ●00		
Dimension of c	classical toric code			

"Recall": The integral points of P give a monomial basis of  $C_P$  and  $PC_P$ .

Integral point 
$$m \in P \cap \mathbb{Z}^N \Leftrightarrow \operatorname{ev}\left(\chi^{(m,P)}\right) \in \mathsf{C}_P/\mathsf{P}\mathsf{C}_P$$

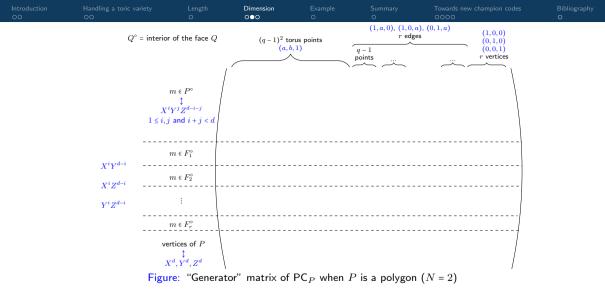
CLASSICAL CASE: on  $\mathbb{F}_q^*$ ,  $x^{q-1} = 1$ . For two elements  $(u, v) \in (\mathbb{Z}^N)^2$ , we write  $u \sim v$  if  $u - v \in (q-1)\mathbb{Z}^N$ .

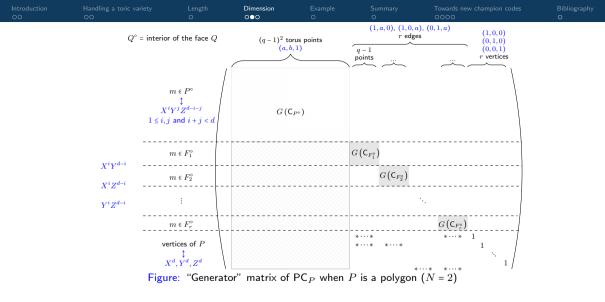
## Theorem [Ruano 07]

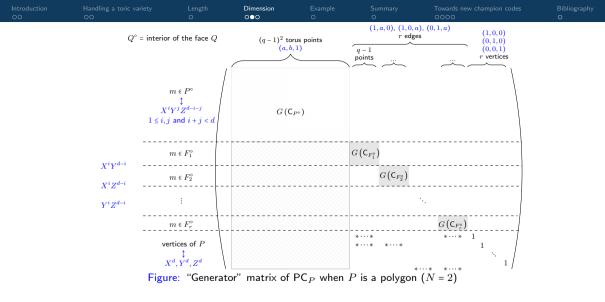
• 
$$\chi^{(m,P)}(\mathbf{t}) = \chi^{(m',P)}(\mathbf{t})$$
 for every  $\mathbf{t} \in \mathbb{T}_P(\mathbb{F}_q) \Leftrightarrow m \sim m'$ ,

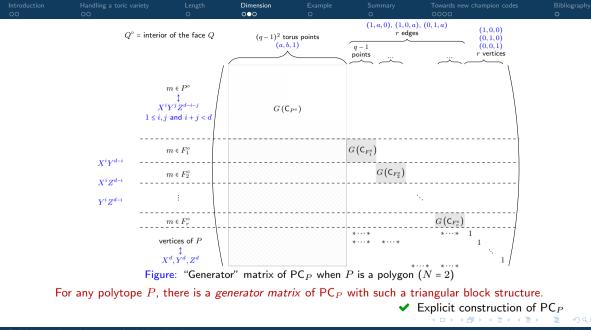
• If  $\overline{P}$  is a set of representatives of  $P \cap \mathbb{Z}^N$  modulo ~, then  $\{(\chi^{\langle \overline{m}, P \rangle}(\mathbf{t}), \mathbf{t} \in \mathbb{T}_P(\mathbb{F}_q) \mid \overline{m} \in \overline{P}\}$  is a basis of  $C_P$ .

Not so nice when homogenizing! On  $\mathbb{P}^1(\mathbb{F}_q)$ ,  $X_0^q \neq X_0 X_1^{q-1}$  at [1:0].









			Dimension			
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Dimension and	d reduction modulo g	l - 1				

Dimension of  $PC_P$  = rank of the previous matrix =  $\sum_Q \dim C_{Q^\circ}$ 

PROJECTIVE CASE: Reduction of P face by face.

On  $P \cap \mathbb{Z}^N$ , we write  $m \sim_P m'$  if there exists a face Q of P s.t.  $m, m' \in Q^\circ$  and  $m - m' \in (q - 1)\mathbb{Z}^N$ .

## Theorem [N. 20]

- $\chi^{\langle m, P \rangle}(\mathbf{x}) = \chi^{\langle m', P \rangle}(\mathbf{x})$  for every  $\mathbf{x} \in \mathbf{X}_P(\mathbb{F}_q) \Leftrightarrow m \sim_P m'$ ,
- If  $\operatorname{Red}(P)$  is a set of representatives of  $P \cap \mathbb{Z}^N$  modulo  $\sim_P$ , then  $\left\{ \operatorname{ev}_P(\chi^{(\overline{m},P)} | \overline{m} \in \operatorname{Red}(P) \right\}$  is a basis of  $\mathsf{PC}_P$ .

✓ Dimension of  $PC_P$ 

Let  $a, b, \eta \in \mathbb{N}^*$  and  $P(\eta) = \operatorname{Conv}((0,0), (a,0), (a,b), (0, b + \eta a)).$  $\rightarrow \mathbf{X}_{P(\eta)}$  called a *Hirzebruch surface* + a divisor of *bidegree* (a,b).

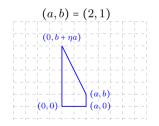
$$\mathbf{X}_{P(\eta)}(\mathbb{F}_q) = (q-1)^2 + 4(q-1) + 4 = (q+1)^2.$$

ightarrow Reduce P modulo q - 1 = 6.

Let us compare the dim  $PC_P$  and dim  $C_P$  on  $\mathbb{F}_7$  for different (a, b).

 $\vdash$  Reduce the interior of each face modulo q - 1 = 6.





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Introduction Handling a toric variety Length Dimension Example Summary Towards new champion codes Bibliography 00 00 00  $\bullet$  000  $\bullet$  0000 0 Example of computation of the dimension of PC<sub>P</sub> and C<sub>P</sub>

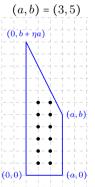
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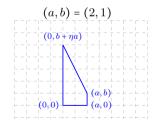
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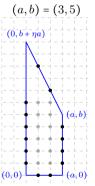
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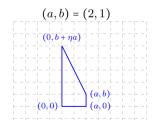
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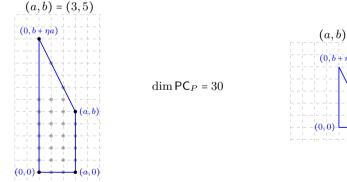
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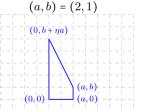
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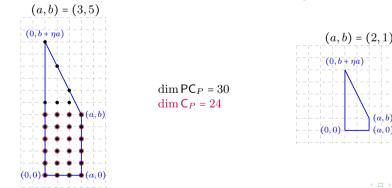


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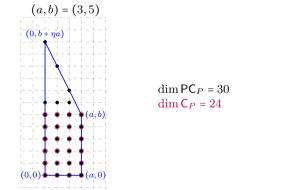


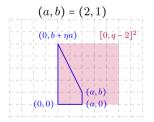
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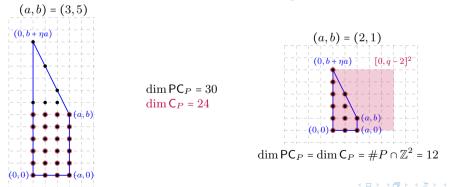


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			Summary •	
Minimum dista	ance			

Lower bound on the minimum distance of  $PC_P$  more technical [CN16, Nar19] *Key ingredient:* (theorical) **Gröbner basis** of the vanishing ideal of  $\mathbf{X}_P(\mathbb{F}_q)$  $\rightarrow$  no problem from the exponential growth in #variables of the complexity of its actual computation.

In conclusion, this work provides a general framework for studying AG codes on toric varieties. Given a polytope P, we can

- compute exactly the dimension of the code PC<sub>P</sub>,
- get a lowerbound on the minimum distance (not always sharp),

provided that we have a good algorithm to determine the integral points of a polytope.

 $\tilde{O}\left(\left(s^{\lceil \frac{N}{2} \rceil} + V\right)\log \delta\right)$  for a polytope of dim. N of vol. V with s vertices, and where  $\delta$  is the maximum modulus of the coordinates of the vertices of P [SV13, Prop. 3.5].

			Towards new champion codes	
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Brown and Kasprzyk [BK13] systematically investigated (generalized) toric codes associated to small polygons  $\rightarrow$  good codes acheiving/beating the best-known parameters.

Given a champion toric code  $C_P$ ,

- $\ominus$  PC<sub>P</sub> is unlikely to be a champion code itself,
- $\oplus$  indicate how to extend C<sub>P</sub> while keeping good parameters.

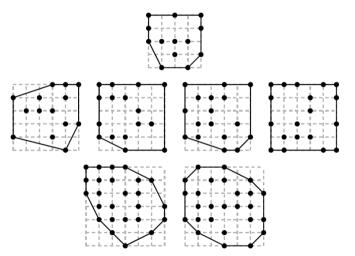
Champion generalizing toric code [49, 14, 26] over  $\mathbb{F}_8$  [BK13] Cannot consider its convex hull (simplicial toric variety on  $\mathbb{F}_8$ )  $\rightarrow$  projective toric code on  $\mathbb{X}_P$  but PC<sub>P</sub> is [87, 14, 34]<sub>8</sub>.

Let us puncture this code!



Figure: A polygon containing the points defining a champion generalized toric code [49, 14, 26] over  $\mathbb{F}_8$  [BK13]

						Towards new champion codes	
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7 ways to ge	et non Hamming-equi	valent (gene	eralized) [49.1	4.26] toric	codes [BK13		



			Towards new champion codes ○○●○	
What now?				

- Looking for new champion codes this way...
- Investigate properties of these codes : Local decodability [LN20],dual codes for application to secret sharing [Han16]

Thank you!

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KEY INGREDIENT: Gröbner basis of the vanishing ideal of  $\mathbf{X}_{P}(\mathbb{F}_{q})$  [CN16, Nar19] • Choose a nice total order < on  $\mathbb{Z}^{N}$  (addition compatibility) : lexicographic

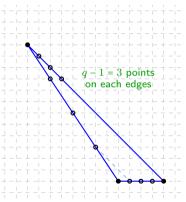
- Find  $\lambda$  s.t. for every face Q of  $\lambda P$ ,  $\# \operatorname{Red}(Q^\circ) = (q-1)^{\dim Q}$ (*i.e.*  $\mathsf{PC}_{\lambda P} = \mathbb{F}_q^n$ )
- Compute Red(P) and Red(λP) taking into account the order.
   Representative = smallest element wrt < among a class</li>

modulo  $\sim_{(\lambda)P}$ 

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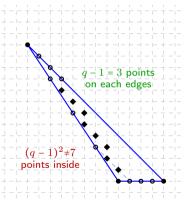
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♥ Compute  $\operatorname{Red}(P)$  and  $\operatorname{Red}(\lambda P)$  taking into account the order.



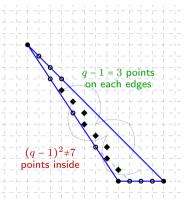
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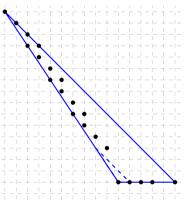
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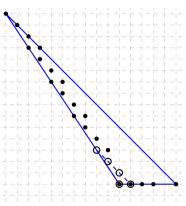
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$$d(\mathsf{PC}_P) \geq \min_{m \in \mathrm{Red}_{<}(P)} \# \left( (m + P_{\mathsf{surj}} - P) \cap \mathrm{Red}_{<}(P_{\mathsf{surj}}) \right).$$



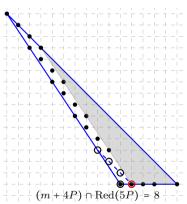
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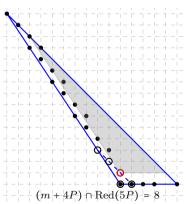
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